



Managing the Endianness of Software Building Blocks with GNAT Ada Attributes: a Case Study

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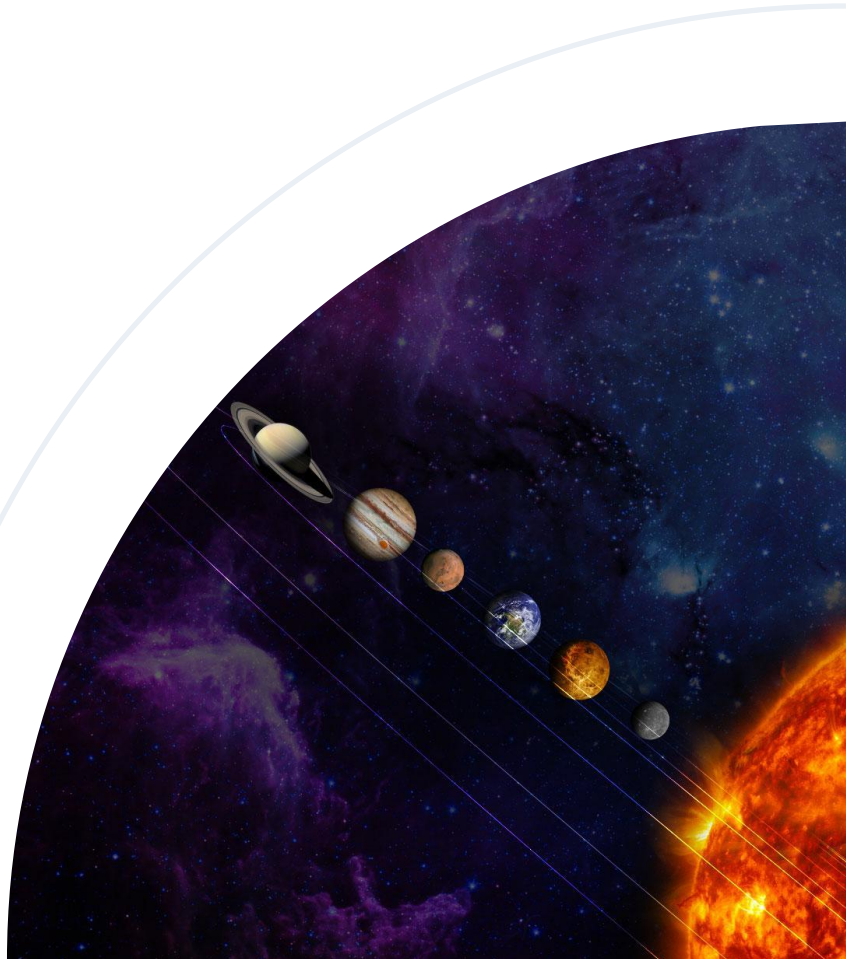
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Outline

- 🪐 Recap on Endianness
- 🪐 Context
- 🪐 The PUS standard
- 🪐 Scalar Storage Order Attribute
- 🪐 Application to selected case study
- 🪐 Conclusions

🪐 (+) We acknowledge Adrien Prioux, now with Thales Services, for his essential contribution to this investigation



Endianness

Big endian

Big-endian format requires storage with the most significant byte first. Namely, that the most significant bit is stored at the lowest memory address then the following bytes are stored in decreasing order of significance.

Example : **one hundred twenty-three**

BIG ENDIAN

123



Little endian

Little-endian format reverses the order of the sequence and stores the least significant byte at the first location in memory and the most significant byte is stored last.

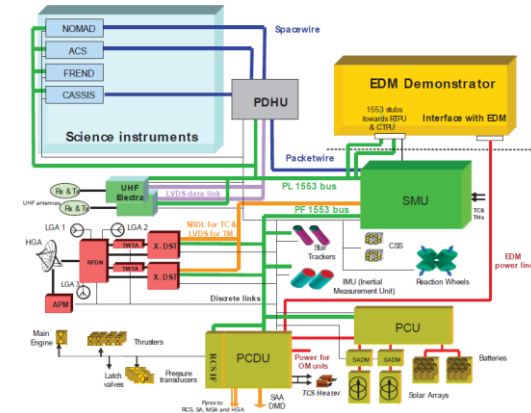
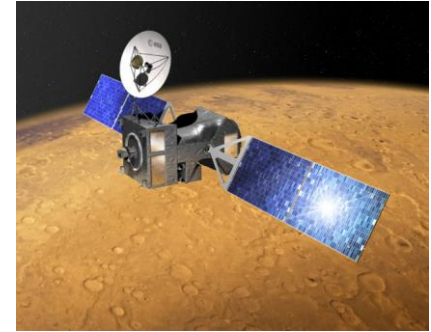
LITTLE ENDIAN

321

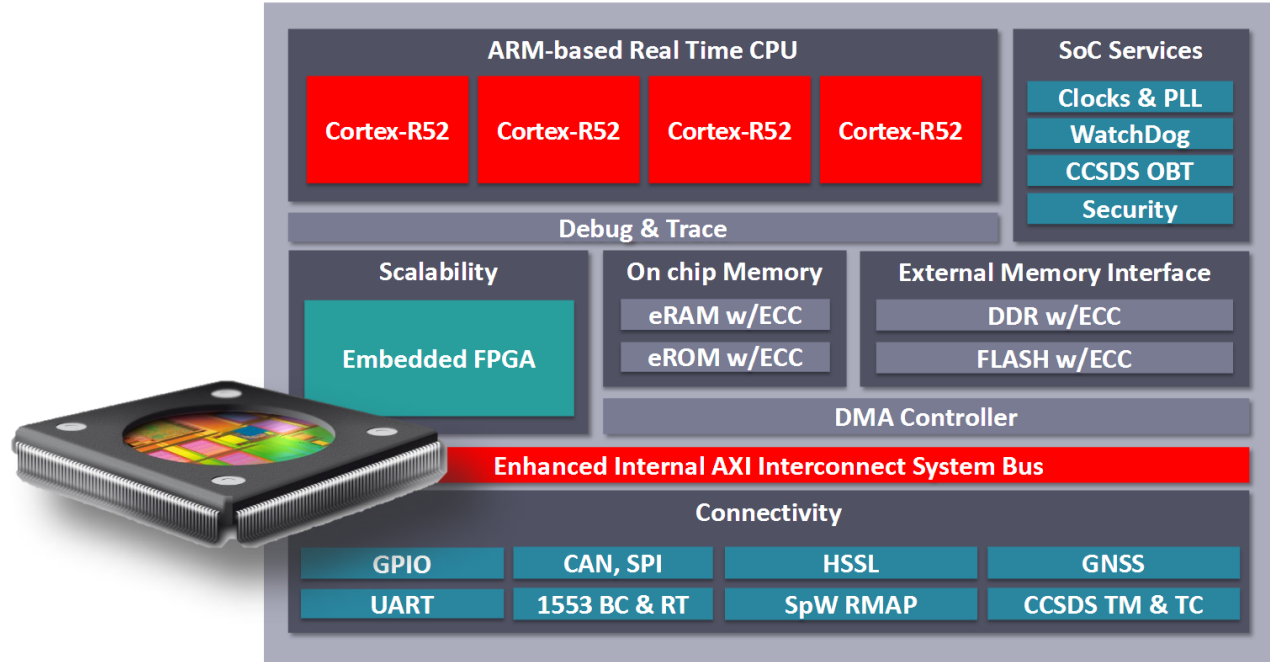


Why is this relevant to space software?

- 🚀 Current US or EU spacecrafts mostly use Big-Endian Processor Architectures
 - 🚀 SPARC (mostly in Europe, some cases in US)
 - 🚀 ERC32, LEON2, LEON3, LEON4
 - 🚀 PowerPC (US, Europe in some cases)
 - 🚀 PPC 750, PPC 8548, etc..
- 🚀 New ARM-based hardware targets (Little Endian) will start to be used in the next future
 - 🚀 ARM R5 (Project developments already ongoing)
 - 🚀 ARM R52 (H2020 DAHLIA)
 - 🚀 ARM A9 (CNES Hyperion / ESA OPSAT)
 - 🚀 ARM A53 (new NASA on-board computer)
- 🚀 Porting consolidated building blocks with 10+ years of heritage across architectures requires to address the problem of endiannes in a satisfactorily manner



Example of next ARM hardware target



Goals and case study

- 🪐 Facilitate retargeting from current big-endian hardware to another platform with little-endian hardware
 - 🪐 Maintain a unique version of OBSW code which works on both architectures
 - 🪐 Produce an endian-agnostic application and solve interoperability with ground software
 - 🪐 Remain compatible with Ada 2005 and with both GNAT compiler for ARM and SPARC architectures
 - 🪐 Come up with a solution should be affordable and easy to implement
-
- 🪐 Selected case study: application to a subset of the TAS **PUS library**
 - 🪐 Overall 30000 Ada SLOCs

PUS

“Telemetry and Telecommand Packet Utilisation”
ECSS-E-ST-70-41C

The European standard to describe the contents and format of spacecraft Telecommands (TC) and Telemetry (TM)

- Description of common services to operate the spacecraft

- Mostly focuses on the data format of the exchange

- With some data exchanges protocol

Packet Header (48 Bits)						Packet Data Field (Variable)					
Packet ID				Packet Sequence Control		Packet Length	Data Field Header (Optional) (see Note 1)	Application Data	Spare	Packet Error Control (see Note 2)	
Version Number (=0)	Type (=1)	Data Field Header Flag	Application Process ID	Sequence Flags	Sequence Count						
3	1	1	11	2	14						
16				16		16	Variable	Variable	Variable	16	



Space engineering
Telemetry and telecommand packet utilization

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service type	
name	ID
request verification	1
device access	2
housekeeping	3
parameter statistics reporting	4
event reporting	5
memory management	6
(reserved)	7
function management	8
time management	9
(reserved)	10
time-based scheduling	11
on-board monitoring	12
large packet transfer	13
real-time forwarding control	14
on-board storage and retrieval	15
(reserved)	16
test	17
on-board control procedure	18
event-action	19
parameter management	20
request sequencing	21
position-based scheduling	22
file management	23

Note: The reserved service type identifiers were used in previous versions of this Standard. This Standard no longer promotes the use of these service types but does not preclude that existing implementations are reused for new missions.

PUS – Datafield examples

8.1.2.8 TM[1,8] failed completion of execution verification report

- a. Each telemetry packet transporting a failed completion of execution verification report shall be of message subtype 8.

NOTE For the corresponding system requirements, refer to clause 6.1.5.3.2.

- b. For each telemetry packet transporting a failed completion of execution verification report, the source data field shall have the structure specified in Figure 8-8.

request ID						failure notice	
packet version number	packet ID			packet sequence control		code	data
	packet type	secondary header flag	application process ID	sequence flags	packet sequence count		
enumerated (3 bits)	enumerated (1 bit)	Boolean (1 bit)	enumerated (11 bits)	enumerated (2 bits)	unsigned integer (14 bits)	enumerated	deduced

deduced presence

NOTE The request ID field alone cannot be used to identify the request since it does not contain the identifier of the source of that request. That source identifier corresponds to the destination identifier of the secondary header of the related telemetry packet, refer to clause 7.4.3.1.

8.22.2.10 TM[22,10] position-based schedule detail report

- a. The telemetry packet transporting a position-based schedule detail report shall be of message subtype 10.

NOTE For the corresponding system requirements, refer to clause 6.22.9.2.

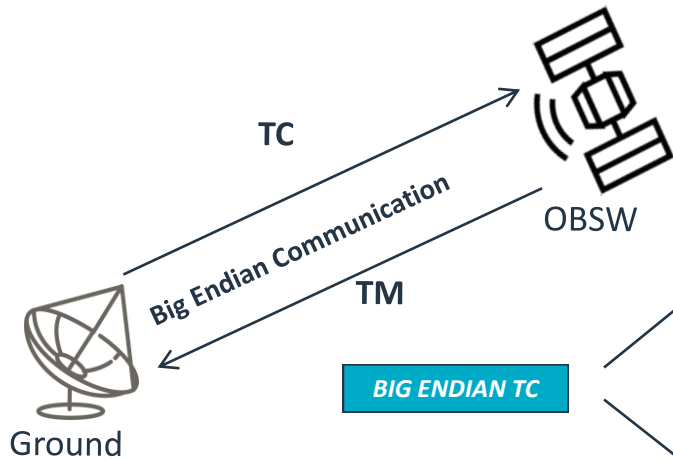
- b. For each telemetry packet transporting a position-based schedule detail report, the source data field shall have the structure specified in Figure 8-246.

<i>repeated N times</i>						
N	sub-schedule ID	group ID	position tag	activity persistency status	persistent activity periodicity	request
enumerated	enumerated	enumerated	deduced	enumerated	unsigned integer	TC packet
					<i>deduced presence</i>	
	<i>optional</i>	<i>optional</i>		<i>optional</i>		

NOTE 1 The structure of the position tag field is driven by requirement 8.22.1b.

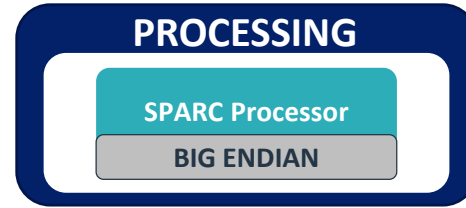
NOTE 2 For the activity persistency status enumerated values, refer to requirement 8.22.3d.

Big Endian vs. Little Endian



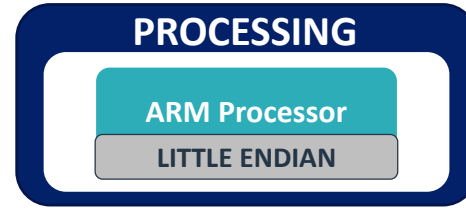
BIG ENDIAN TC

Example : **0x12345678**



VALID DATA

Big-Endian	1 Byte	1 Byte	1 Byte	1 Byte
Memory Address	0	1	2	3
32-bit integer	12	34	56	78
Read-Value	0x12345678			



WRONG DATA

Little-Endian	1 Byte	1 Byte	1 Byte	1 Byte
Memory Address	0	1	2	3
32-bit integer	12	34	56	78
Read-Value	0x78563412			

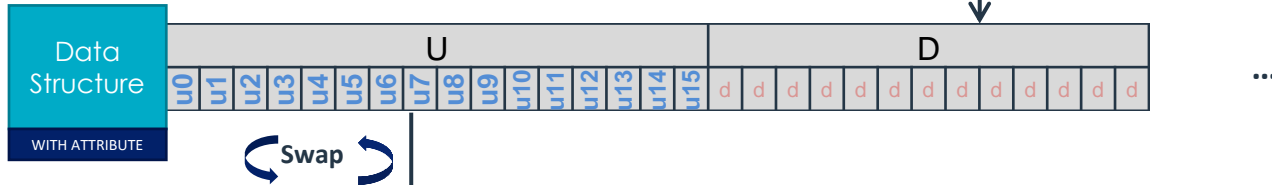
TC reception

Received data – Big Endian (Byte-stream received from low-level drivers)



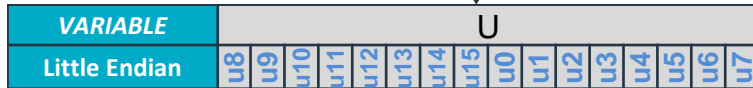
Unchecked conversion to endian pragma structure

Unchecked conversion



Assign to a variable

Assign an element to a variable



TC to Data Structure

```
'type TC_HEADER_T is record
  U : UINT16_T;
  D :  UINT13_T;
  T :  UINT5_T;
  Q :  UINT6_T;
  C :  UINT8_T;
end record;

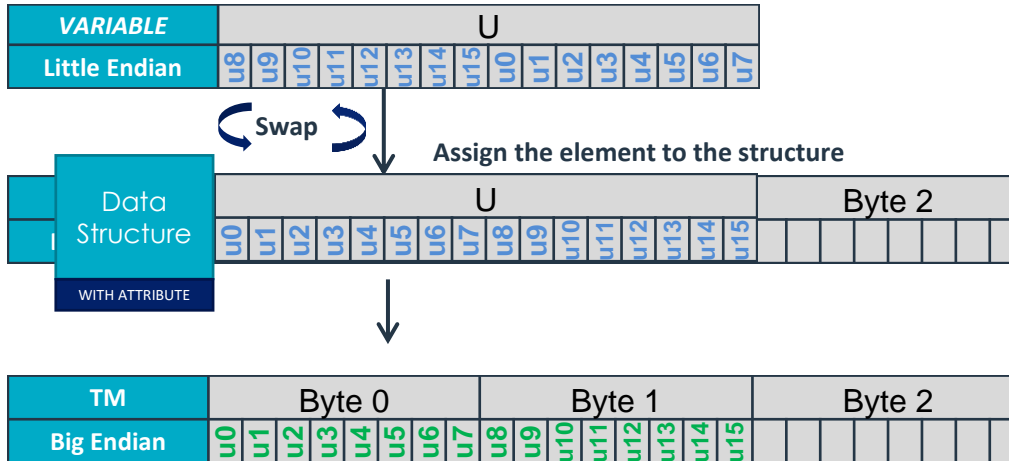
for TC_HEADER_T use record
  U at 0 range 0 .. 15;
  D at 0 range 16 .. 28;
  T at 0 range 29 .. 33;
  Q at 0 range 34 .. 39;
  C at 0 range 40 .. 47;
end record;

for TC_HEADER_T'Bit_Order use System.High_Order_First;
for TC_HEADER_T'Scalar_Storage_Order use System.High_Order_First;
```

```
--Received data Exemple component U is 16#1234#
subtype TC_T is UINT8_ARRAY_NC_T(0 .. 5);
TC      : TC_T := (16#12#,16#34#,16#56#,16#78#,16#9a#,16#bc#);
--Header Structure
TC_HEADER : TC_HEADER_T;
-- to convert the TC packet from bytes list to an TC_HEADER_T structure
function TO_TC_HEADER is new UNCHECKED_CONVERSION (TC_T, TC_HEADER_T);
--Memory map of TC_HEADER : 16#12# , 16#34#
TC_HEADER := TO_TC_HEADER (TC);
--Memory map of U_var : 16#34# , 16#12#
U_var := TC_HEADER.U;
```

```
<gdb> x /2xb &TC_HEADER
0x245fe78:      0x12      0x34
<gdb> x /2xb &U_var
0x245fe8a:      0x34      0x12
```

On-board parameter sent as TM



Assembler level

	BIG ENDIAN	LITTLE ENDIAN
X86	<pre>401a8e: 8b 45 ec mov -0x14(%ebp),%eax 401a91: 89 45 e0 mov %eax,-0x20(%ebp)</pre>	<pre>401a8e: 8b 45 ec mov -0x14(%ebp),%eax 401a91: 0f c8 bswap %eax 401a93: 89 45 e0 mov %eax,-0x20(%ebp)</pre>
ARM	<pre>3b48: 689b ldr r3, [r3, #8] 3b4a: 461a mov r2, r3</pre>	<pre>3b48: 689b ldr r3, [r3, #8] 3b4a: ba1b rev r3, r3 3b4c: 461a mov r2, r3</pre>

```
NB_TC_SELECTION_REC : NB_TC_SELECTION_REC_T
:= ( NUMBER_OF_TC , 16#2233#, 16#4455# );
```

```
401a10: 66 c7 45 e2 33 22  movw $0x2233,-0x1e(%ebp)
401a16: c7 45 e4 55 44 00 00  movl $0x4455,-0x1c(%ebp)
```

```
401a10: 66 c7 45 e2 22 33  movw $0x3322,-0x1e(%ebp)
401a16: c7 45 e4 00 00 44 55  movl $0x55440000,-0x1c(%ebp)
```

“Propagation” of attribute specification

```
type TC_HEADER_T is
record
  CCSDS_HEADER : CCSDS_TC_HEADER_T;
  PUS_HEADER   : PUS_TC_HEADER_T;
end record;

for TC_HEADER_T use
record
  CCSDS_HEADER at 0 range 0 .. 47;
  PUS_HEADER   at 0 range 48 .. 71+8*PUS.CST.SRC_DEST_ID_FIELD_SIZE_C;
end record;

for TC_HEADER_T SIZE use 72+8*PUS.CST.SRC_DEST_ID_FIELD_SIZE_C;
for TC_HEADER_T Bit_Order use System.High_Order_First;
for TC_HEADER_T Scalar_Storage_Order use System.High_Order_First;
```

```
type CCSDS_TC_HEADER_T is
record
  PACKET_ID       : CCSDS_TC_PACKET_ID_T;
  SEQUENCE_FLAGS : PUS.TYPES.BITS2_T;
  SEQUENCE_COUNT : PUS.TYPES.SEQUENCE_COUNT_T;
  PACKET_LENGTH  : PUS.TYPES.UNSIGNED_HALF_WORD_T;
end record;

for CCSDS_TC_HEADER_T use
record
  PACKET_ID       at 0 range 0 .. 15;
  SEQUENCE_FLAGS at 0 range 16 .. 17;
  SEQUENCE_COUNT at 0 range 18 .. 31;
  PACKET_LENGTH  at 0 range 32 .. 47;
end record;

for CCSDS_TC_HEADER_T SIZE use 48;

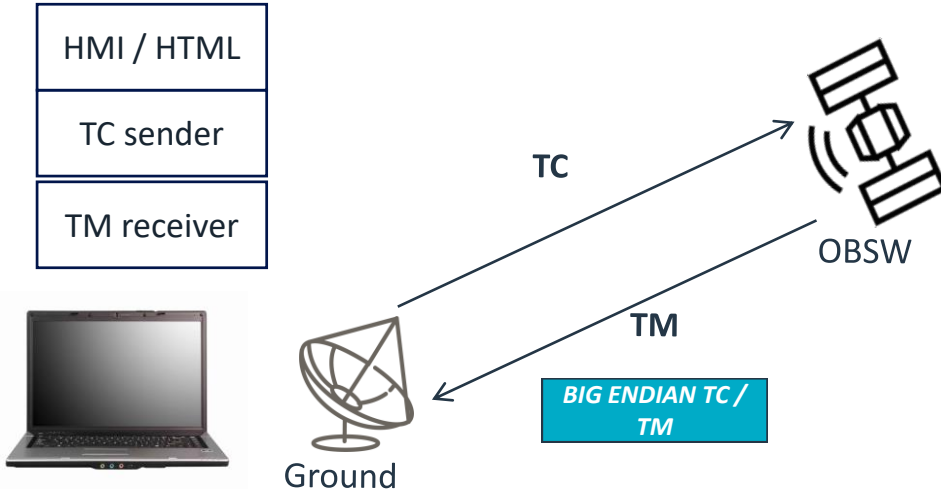
for CCSDS_TC_HEADER_T Bit_Order use System.High_Order_First;
for CCSDS_TC_HEADER_T Scalar_Storage_Order use System.High_Order_First;
```

```
type PUS_TC_HEADER_T is
record
  SECONDARY_HEADER_FLAG : PUS.TYPES.BIT_T;
  VERSION               : PUS.TYPES.BITS3_T;
  ACK_COMPLETE         : PUS.TYPES.BIT_T;
  ACK_PROGRESS         : PUS.TYPES.BIT_T;
  ACK_START            : PUS.TYPES.BIT_T;
  ACK_ACCEPTANCE      : PUS.TYPES.BIT_T;
  SERVICE_TYPE        : PUS.TYPES.SERVICE_TYPE_T;
  SERVICE_SUBTYPE     : PUS.TYPES.SERVICE_SUBTYPE_T;
  SRC_DEST_FIELD_USED  -- if source id/destination id fields are included in the PUS header
  SOURCE_ID            : PUS.TYPES.BYTE_T;
end if;
end record;

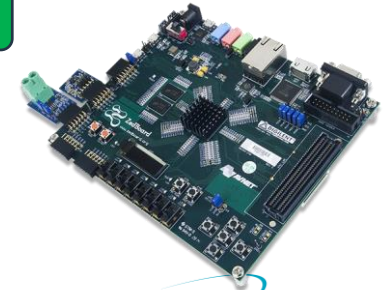
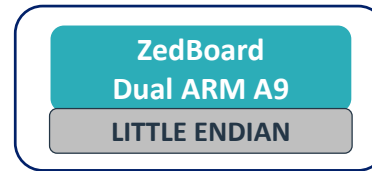
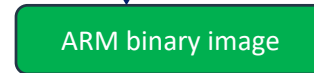
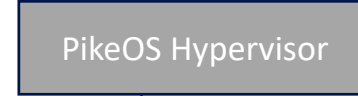
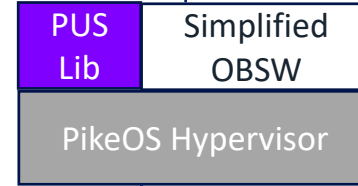
for PUS_TC_HEADER_T use
record
  SECONDARY_HEADER_FLAG at 0 range 0 .. 0;
  VERSION               at 0 range 1 .. 3;
  ACK_COMPLETE         at 0 range 4 .. 4;
  ACK_PROGRESS         at 0 range 5 .. 5;
  ACK_START            at 0 range 6 .. 6;
  ACK_ACCEPTANCE      at 0 range 7 .. 7;
  SERVICE_TYPE        at 0 range 8 .. 15;
  SERVICE_SUBTYPE     at 0 range 16 .. 23;
  SRC_DEST_FIELD_USED  -- if source id/destination id fields are included in the PUS header
  SOURCE_ID            at 0 range 24 .. 31;
end if;
end record;

for PUS_TC_HEADER_T SIZE use 8*PUS.CST.PUS_HEADER_SIZE_C;
for PUS_TC_HEADER_T Bit_Order use System.High_Order_First;
for PUS_TC_HEADER_T Scalar_Storage_Order use System.High_Order_First;
```


Case study execution



Tested on a subset of PUS services:
1,3, 5, 17



Conclusions

- 🌐 This study had the goal to investigate new techniques for managing endianness of code
 - 🌐 By maintaining a single code baseline for both big-endian and little-endian architectures
 - 🌐 With a solution that could be as simple and elegant as possible
- 🌐 We therefore focused on the attribute “Scalar Storage Order” attribute
- 🌐 The prototype realised during the investigation permitted to confirm the correct behavior of the attribute
- 🌐 The solution is more interesting than some others (e.g., macros or code-swap functions)
 - 🌐 Applies to data declaration
 - 🌐 A priori does not impact functional / business code
- 🌐 However a toll has to be paid
 - 🌐 In some cases, ad-hoc intermediate data types have to be introduced in order to “force endianness checks”
 - 🌐 E.g., to cope with extraction of scalar values with a given endianness
- 🌐 The solution works well if the overall software architecture was already layered so as to isolate most of the endianness concerns in 1-2 modules of the architecture

End of Presentation

QUESTIONS?



OPEN